### An interaction between queing and change detection

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## THE DYNAMICS of a QUEUE

Consider a single server queue. Waiting time of the *n*-th customer:  $W_n$ . The dynamics of  $W_n$  is given by a *non-linear system*:

 $W_n = (W_{n-1} + X_n)_+$  with  $W_0 = 0,$  (1)

where  $X_n = V_{n-1} - U_n$  = service time minus interarrival time.

A system-theoretic point of view: (8) is not a stable system.

A similar non-linear dynamics arises in the theory of risk processes:

$$W_n^- = (W_{n-1}^- + X_n^-)_-$$
 with  $W_0^- = K > 0.$  (2)

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# PROBABILISTIC STABILITY of a QUEUE

A standard assumption: assume i.i.d. inputs  $(X_n)$ , with  $E(X_n) < 0$ .

Markovian techniques: establish geometric ergodicity assuming

$$\mathbb{E}(\exp c' X_1) < 1 \quad \text{for some} \quad c' > 0. \tag{3}$$

Strong LLN follows for functions of  $W_n$ . See Meyn & Tweedie.

# PROBLEM STATEMENT

Under what conditions for the inputs  $(X_n)$  can we ensure:

1. A strong LLN for the empirical tail probabilities:

$$\limsup_{N} \frac{1}{N} \sum_{n=1}^{N} \mathbb{I}_{\{W_n > K\}} \leq \limsup_{n} P(W_n > K) \quad \text{a.s.}$$

2. Exponential decay of tail probabilities:

 $P(W_n > K) < Ce^{-cK}.$ 

Technical kinship of the two problems.

The problem: detect changes of statistical patterns of signals in real time.

Example: monitoring EEG signals for epileptic patients

See: V.Poor and O.Hadjiliadis (2009): Quickest Detection.

# A CLASSIC PROBLEM

Given a sequence of i.i.d. r.v.-s  $Y_n$  with prob. density functions

 $f(y, \theta_0)$  for  $n < \tau$  and  $f(y, \theta_1)$  for  $n \ge \tau$ .

Estimate

change point : au

using observations  $(y_n)$ .

The Cumulative Sum (CUSUM) test or Page-Hinkley detector:

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E.S. Page, Biometrika, 1954
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D.V. Hinkley, J.Amer. Statist. Assoc., 1971

# STATISTICS and IT

A modern interpretation, following J.Rissanen, 1989: Encode data using the two possible models, following Inf.Thy.: The quasi-optimal code-lengths are

 $-\log f(y_n, \theta_0)$  and  $-\log f(y_n, \theta_1)$ .

The differences in code-lengths define the score

 $X_n = -\log f(y_n, \theta_0) + \log f(y_n, \theta_1).$ 

Now the information inequality gives

$$\mathbf{E}X_n < \mathbf{0}$$
 for  $n < \tau$  and  $\mathbf{E}X_n > \mathbf{0}$  for  $n \ge \tau$ .

# THE CUSUM TEST for I.I.D. DATA

Let  $S_0 = 0$  and let

$$S_n = \sum_{k=1}^n X_k$$
 for  $n \ge 1$ .

Then  $ES_n$  has a minimum at  $\tau - 1$ .

Task: approximate on-line minimization of  $S_n$ .

The CUSUM statistics or Page-Hinkley detector: define

$$g_n=S_n-\min_{0\leq k\leq n}S_k.$$

Generate an alarm if  $g_n > \delta$ , with some fixed threshold  $\delta > 0$ .

## FALSE ALARM RATE

Apply the Page-Hinkley detector to a process with no change at all. A key performance characteristics: false alarm probability

 $\limsup_n P_{\theta_0}(g_n > \delta).$ 

Practical relevance: false alarm rate (FAR) defined as a.s.

$$\limsup_{N} \frac{1}{N} \sum_{n=1}^{N} \mathbb{I}_{\{g_n > \delta\}}.$$

Problem: find an upper bound for the FAR.

# THE DYNAMICS of CUSUM

The dynamics of  $g_n$  is easily obtained as follows:

$$g_n = (g_{n-1} + X_n)_+$$
 with  $g_0 = 0$ .

This establishes the link between queuing,  $W_n$ , and change detection,  $g_n$ .

Objective: bounding the empirical tail probabilities of  $g_n$ .

Two related technical problems:

- Exponential bounds for the tail probabilities of g<sub>n</sub>.
- Mixing properties of g<sub>n</sub>.

## MOTIVATION for MIXING

Let  $(\nu_n)$  be an  $\mathbb{R}^s$ -valued i.i.d. sequence of r.v.-s such that

$$\sup_{n\geq 0} \mathbb{E} |\nu_n|^q < +\infty \quad \text{ for all } \quad 1\leq q<\infty.$$

Let the  $s \times s$  matrix A be stable, and define the filtered process

$$X_n = AX_{n-1} + \nu_n \quad \text{with} \quad X_0 = 0.$$

Decompose  $(X_n)$  as

$$X_n = A^{\tau} X_{n-\tau} + \sum_{k=0}^{\tau-1} A^k \nu_{n-k}.$$

# L-MIXING, I.

#### Definition

Let  $X = (X_n)$  be a stochastic process on  $(\Omega, \mathcal{F}, \mathbb{P})$ . X is *M*-bounded if for all  $1 \leq q < +\infty$ 

$$M_q(X) := \sup_{n\geq 0} \|X_n\|_q < +\infty.$$

Let  $\nu = (\nu_n)$  be an i.i.d. sequence, and define its past and future as

 $\mathfrak{F}_n = \sigma(\nu_k : k \leq n) \text{ and } \mathfrak{F}_n^+ = \sigma(\nu_k : k \geq n+1).$ 

Let  $\tau > 0$  be an integer, a fixed memory length, and defined for  $1 \le q < +\infty$  the error of approximation by the near past as

$$\gamma_q(\tau, X) = \gamma_q(\tau) := \sup_{n \ge \tau} \| X_n - \operatorname{E}(X_n | \mathcal{F}_{n-\tau}^+) \|_q.$$

# L-MIXING, II.

#### Definition

A stochastic process  $X = (X_n)$  is *L*-mixing w.r.t.  $(\mathcal{F}_n, \mathcal{F}_n^+)$  if it is adapted i.e.  $X_n$  is  $\mathcal{F}_n$ -measurable for all  $n \ge 1$ , X is *M*-bounded, and

$$\Gamma_q(X) := \sum_{\tau=0}^{+\infty} \gamma_q(\tau) < +\infty \quad \text{for all} \quad 1 \le q < +\infty.$$

Remark: we can also have  $q = \infty$ .

References:

Ljung, L.: Math. Programming, 1976.

LG: Stochastics, 1989.

# THM: for I.I.D. SCORES $g_n$ is L-MIXING

Assume  $\mathbb{E}(X_1) < 0$ , and also  $\mathbb{E}(\exp cX_1) < \infty$  for some c > 0. Then

 $\mu := \mathbb{E}(\exp c' X_1) < 1 \quad \text{for some} \quad c' > 0. \tag{4}$ 

Let  $\mathfrak{F}_n := \sigma(X_i \mid i \leq n)$  and  $\mathfrak{F}_n^+ := \sigma(X_i \mid i \geq n+1)$ .

#### Theorem

Let  $(X_n)$  be a sequence of i.i.d. random variables such that (4) holds. Then the Page-Hinkley detector  $(g_n)$  is L-mixing with respect to  $(\mathcal{F}_n, \mathcal{F}_n^+)$ .

LG and Prosdocimi, C. Systems & Control Letters, 2011.

## EMPIRICAL TAIL PROBABILITIES

A known result in risk theory: for any c'' such that 0 < c'' < c', we have

$$\sup_{n} \mathbb{E}\left(\exp c'' g_{n}\right) < \infty.$$
(5)

Hence  $P(g_n > \delta) \le Ce^{-c''\delta}$  with some  $0 < C < \infty$ . Equivalently,

 $\mathbb{E} \mathbb{I}_{\{g_n > \delta\}} \leq C e^{-c''\delta}.$ 

Since  $(g_n)$  is *L*-mixing, by a strong LLN it follows that

$$\limsup_{N} \frac{1}{N} \sum_{n=1}^{N} \mathbb{I}_{\{g_n > \delta\}} \leq C' e^{-c'' \delta}.$$

# DEPENDENT SCORES

The key technical condition to be established above: for any c' > 0

$$\mathbb{E}\exp\sum_{k=1}^n c' X_k \leq C e^{-c'' n}$$

with some C, c'' > 0. Writing the l.h.s. as

$$\mathbb{E}\exp\left(\sum_{k=1}^n c'(X_k - \mathbb{E}X_k)\right) \cdot \exp\left(\sum_{k=1}^n c' \mathbb{E}X_k\right),$$

assuming  $\mathbb{E}X_k \leq -\varepsilon$ , it is sufficient to show that for small c'-s and all n:

$$\mathbb{E}\exp\left(\sum_{k=1}^n c'(X_k - \mathbb{E}X_k)\right) \le e^{\kappa(c')^2 n},$$

with some  $\kappa > 0$ .

# AN EXPONENTIAL MOMENT CONDITION

Let  $X = (X_n)$  be a sequence of real-valued random variables.

**Condition E.** There exist c > 0 and  $\kappa > 0$  such that for  $0 \le c' < c$  and all  $1 \le m \le n$  we have

$$\mathbb{E} \exp\left(c'\sum_{k=m}^{n}(X_k-\mathbb{E}X_k)\right) \leq \exp\left(\kappa \ (c')^2(n-m+1)\right).$$

Objective: find useful conditions for the above inequality to hold.

# I.I.D. REVISITED

#### Lemma

Let  $(X_n)$  be a zero-mean, i.i.d. sequence such that  $\mathbb{E} e^{c|X_n|} < \infty$ . Then Condition E is satisfied.

The proof is trivial, noting:

 $\mathbb{E}e^{c'X} \leq e^{\kappa(c')^2}.$ 

for  $|c'| \leq c$  with some  $\kappa > 0$ .

We get even more.

# ANOTHER EXPONENTIAL MOMENT INEQUALITY

Let X be a two-sided i.i.d. sequence as above. Let  $h = (h_k), k = 0, 1, ...$  be an  $l_1$  - sequence and define

$$Y_n = \sum_{k=0}^{\infty} h_k X_{n-k}$$
 in short  $Y = h \star X$ .

Write 
$$|| h ||_2^2 = \sum_{k=0}^{\infty} h_k^2$$
.

#### Theorem

Let  $Y = (Y_n)$  be a as above. Then for  $||h||_2 \le c$ 

 $\mathbb{E}\exp\left(h\star X\right) \leq \exp\left(\kappa \parallel h \parallel_2^2\right).$ 

# A SECOND EXPONENTIAL MOMENT CONDITION

Let  $X = (X_n)$  be a two-sided sequence of real-valued r.v.-s,  $\mathbb{E}X_n = 0$ .

**Condition SE.** There exist c > 0 and  $\kappa > 0$  such that for any  $h \in I_1$  with  $\|h\|_2 \le c$  we have

$$\mathbb{E}\exp\left(\sum_{k=0}^{\infty}h_{k}X_{n-k}\right)\leq\exp\left(\kappa\parallel h\parallel_{2}^{2}\right).$$
(6)

Note: if  $X = (X_n)$  satisfies Condition SE and  $g \in I_1$  then

 $Y = g \star X$ 

also satisfies Condition SE. (Trivial). Example:  $|X_n| \leq K$ .

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# A NON-CONSTRUCTIVE EXAMPLE

#### Theorem

Let  $(X_n)$  be a zero-mean L-mixing process such that we have

 $M_{\infty}(X) < +\infty$  and  $\Gamma_{\infty}(X) < +\infty$ .

(7)

Then for any deterministic sequence  $f_n$ 

$$\mathbb{E}\exp\left(\sum_{k=1}^n f_n X_n - 2M_{\infty}(X)\Gamma_{\infty}(X)\sum_{k=1}^n f_n^2\right) \leq 1.$$

It follows that Condition SE is satisfied.

LG: Stochastics and Stochastic Reports, 1991.

# AN OPEN PROBLEM: SERIAL INTERCONNECTION

Let  $X = (X_n)$  be a two-sided sequence of real-valued r.v.-s, satisfying  $\mathbb{E}X_n \leq -\epsilon < 0$  and in addition Condition SE.

Consider a single server queue with waiting time denoted by  $W_n$ .

Let the dynamics of  $W_n$  is given by the non-linear system:

$$W_n = (W_{n-1} + X_n)_+$$
 with  $W_0 = 0.$  (8)

Under what conditions does  $(W_n)$  satisfy Condition SE?

## RECREATIONAL MATH

Establish stability properties of the Page-Hinkley-detector for *deterministic inputs*: assume that  $(X_n)$  is a deterministic sequence satisfying

$$\limsup_{N \longrightarrow +\infty} \frac{1}{N} \sum_{n=1}^{N} X_n < 0.$$
(9)

Let  $(g_n)$  be the response of the Page-Hinkley-detector driven by  $(X_n)$ . Does it follow that

$$\limsup_{N \to \infty} \frac{1}{N} \sum_{n=1}^{N} g_n < \infty?$$
 (10)

# THANK YOU for YOUR ATTENTION !